

A Note on Quasi-Robust Cycle Bases*

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Abstract

We investigate here some aspects of cycle bases of undirected graphs that allow the iterative construction of all elementary elementary cycles. We introduce the concept of quasi-robust bases that generalize the notion of robust bases and demonstrate that a certain class of bases of complete bipartite graph $K_{m,n}$ with $m, n \geq 5$ is quasi-robust but not robust. We furthermore disprove a conjecture for cycle bases of Cartesian product graphs.

Keywords: Cycle space, Cycle basis, robust, quasi-robust, Kainen's Basis, elementary cycle, complete bipartite, Cartesian product

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1 Introduction

Cycle bases [2] are not only an interesting characterization of the structure of graphs by themselves but also provide a basis for computational assessments of the cycle structure of a graph. “Cycle-space algorithms”, for instance, attempt to construct the set of all elementary cycles of a graph from a cycle basis \mathcal{B} by iteratively computing the symmetric difference of an elementary cycle and a basis cycle, subsequently retaining the result if and only if it is again an elementary cycle. If a cycle basis is robust, this approach is successful [5, 9, 1, 11].

Unfortunately, however, very little is known about robust cycle bases beyond a few very special graph classes: As shown in [5], the boundaries of the faces of an embedded planar graph form a robust cycle basis. Corresponding cycle-space algorithms are given in [14, 5]. Furthermore, complete graphs have robust bases that are easy to construct explicitly [9]. Robust bases for a small class of cubic graphs are constructed in [11].

There is, at present, neither an efficient algorithm to construct a robust cycle basis for a given input graph, nor is it known whether robust bases always exist. A major obstacle for the investigation of robust cycle bases is the apparent lack of relationships with other classes of cycle bases that have been explored in much more detail in the past [7, 12]. For instance, Dixon and Goodman [4] conjectured that every strictly fundamental cycle basis is cyclically robust. A counterexample, however, was given in [13]. A more systematic search for connections [11] showed that robust and fundamental cycle bases are also unrelated.

In this contribution we first disprove two conjectures on robust cycle bases. We then introduce a relaxed notion of robustness that is still sufficient for the construction of efficient cycle space algorithms and show that “quasi-robust” cycle bases can be constructed for complete bipartite graphs.

2 Preliminaries

Throughout this contribution, let $G = (V, E)$ be a finite undirected simple 2-connected graph. A (*generalized*) *cycle* in G is an Eulerian subgraph of G , i.e., a subgraph of G in which the degree of every vertex is even. A connected Eulerian subgraph in which every vertex has degree 2 will be called an *elementary cycle*. For simplicity, we identify a subset $E' \subseteq E$ of edges of G with the subgraph $G(E') := (\bigcup_{e \in E'} e, E')$ of G that it defines. In particular, we identify cycles with their edge sets. The symmetric difference of two edge sets E' and E'' will be denoted by $E' \oplus E''$, i.e., we put $E' \oplus E'' := (E' \cup E'') \setminus (E' \cap E'')$. It will sometimes be convenient to identify a cycle by the sequence of vertices traversed in one of the two orientations. We write

$$C = \{x_1x_2, x_2x_3, \dots, x_{k-1}x_k, x_kx_1\} =: [x_1, x_2, \dots, x_k] \quad (1)$$

and use the shorthand $xy = \{x, y\}$ to denote edges as pairs of adjacent vertices.

The power set $\mathfrak{P}(E)$ can be regarded as a vector space over $\text{GF}(2) = \{0, 1\}$ with vector addition \oplus and the trivial multiplication operator $1 \cdot D = D, 0 \cdot D = \emptyset$. The *cycle space* $\mathfrak{C}(G)$ is the subspace of $(\mathfrak{P}(E), \oplus, \cdot)$ that consists of the cycles of G (including the “empty cycle” \emptyset), see e.g. [2]. As every 2-connected graph G is connected, the dimension $\dim_{\text{GF}(2)} \mathfrak{C}(G)$ of its cycle space coincides with its cyclomatic number $\mu(G) := |E| - |V| + 1$, see e.g. [6].

A basis \mathcal{B} of $\mathfrak{C}(G)$ that consists of elementary cycles only is a *cycle basis* of G . For every cycle C , there is a unique subset $\mathcal{B}_C \subseteq \mathcal{B}$ of elementary cycles in \mathcal{B} such that $C = \bigoplus_{C' \in \mathcal{B}_C} C'$ holds.

3 Robust and Quasi-Robust Bases

In the following we build on the discussion in [11], which is based on [9] but uses a somewhat different terminology.

Definition 1. A sequence $\mathcal{S} = (C_1, C_2, \dots, C_k)$ of (not necessarily pairwise distinct) elementary cycles is *well-arranged* if, for each $j \leq k$, the partial sum $Q_j = \bigoplus_{i=1}^j C_i$ is an elementary cycle. The sequence \mathcal{S} is *strictly well-arranged* if, for $2 \leq j \leq k$, $C_j \cap Q_{j-1}$ is a path.

By construction, strictly well-arranged cycle sequences are well-arranged, while the converse is not true [9, 11]. A small example of a sequence \mathcal{S} which is well-arranged but not strictly is shown in Figure 1.

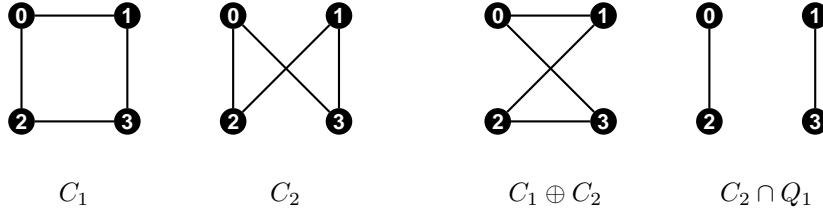


Figure 1: Sequence $\mathcal{S} = (C_1, C_2)$ which is well-arranged, but not strictly well-arranged.

Definition 2. A cycle basis \mathcal{B} is (*strictly*) *quasi-robust* if, for each elementary cycle $C \in \mathcal{C}(G)$ there is a (*strictly*) well-arranged sequence of cycles $\mathcal{S}_C = (C_1, C_2, \dots, C_{k_C})$ with $C_i \in \mathcal{B}$ and $C_{k_C} = C$. The basis \mathcal{B} is (*strictly*) *robust* if for each elementary cycle C , the sequence \mathcal{S}_C can be chosen so that the elementary cycles in \mathcal{S}_C are pairwise disjoint.

Given a graph G , we can associate cycle basis \mathcal{B} with an undirected graph $\Gamma_{\mathcal{B}}$ whose vertices are the elementary cycles in G (including the empty elementary cycle \emptyset). An edge in $\Gamma_{\mathcal{B}}$ connects two elementary cycles C' and C'' if and only if $C' \oplus C'' \in \mathcal{B}$.

Lemma 3. Let \mathcal{B} be a cycle basis of G . Then

1. \mathcal{B} is quasi-robust if and only if $\Gamma_{\mathcal{B}}$ is connected.
2. \mathcal{B} is robust if and only if for every elementary cycle C , the length of a shortest path connecting C with \emptyset equals $|\mathcal{B}_C|$.

Proof. (1) Clear. (2) The path length cannot be smaller than the number $|\mathcal{B}_C|$ of basis cycles necessary to represent C . If C can be reached via no more than $|\mathcal{B}_C|$ intermediate cycles, these must be reached via a well-ordering of \mathcal{B}_C because each basis cycle in \mathcal{B}_C must be used at least once. \square

4 Counterexamples

4.1 Kainen's Basis of $K_{m,n}$

Let us denote by $V_1 \dot{\cup} V_2$ the vertex bipartition of the complete bipartite graph $K_{m,n}$. We fix two vertices $p \in V_1$ and $q \in V_2$ and consider the set of quadrangles

$$\mathcal{B}^{p,q} = \{ \{pq, py, qx, xy\} \mid x \in V_1, y \in V_2 \}. \quad (2)$$

Since the edge xy appears only in a single quadrangle, we see immediately that $\mathcal{B}^{p,q}$ is linearly independent. Furthermore, $|\mathcal{B}^{p,q}| = (|V_1| - 1) \times (|V_2| - 1) = \mu(K_{m,n})$, i.e., $\mathcal{B}^{p,q}$ is basis of $\mathcal{C}(K_{m,n})$, which we will refer to as Kainen's basis.

In [9], it was argued that $\mathcal{B}^{p,q}$ is robust. Here we give a counterexample. In $K_{5,5}$, consider the cycle C as shown in Fig. 2. As indicated in the caption of Fig. 2, its generating set $\mathcal{B}_C^{1,8}$ w.r.t. the Kainen basis $\mathcal{B}^{1,8}$ cannot be well-arranged, hence $\mathcal{B}^{1,8}$ is not robust. We can make an even stronger statement: None of the Kainen bases of $K_{5,5}$ are robust. Choose an arbitrary pair of vertices $p \in V_1$ and $q \in V_2$. Then there is an automorphism π of $K_{5,5}$ with $\pi(p) = 1$ and $\pi(q) = 8$. More detailed informations about graph automorphism can be seen e.g. in [3]. The pre-image of C , $\pi^{-1}(C)$, is necessarily again an elementary cycle. The relationship of $\pi^{-1}(C)$ and $\mathcal{B}^{p,q}$ is the same as that of C and $\mathcal{B}^{1,8}$, hence the generating set of $\pi^{-1}(C)$ cannot be well-arranged, implying that $\mathcal{B}^{p,q}$ is also not robust.

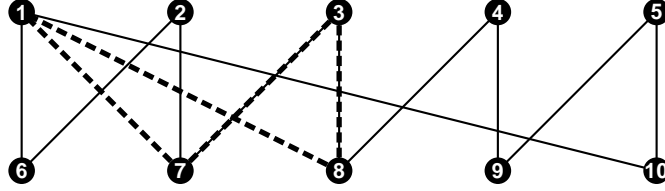


Figure 2: Counterexample for Kainen's assertion. Consider the basis $\mathcal{B}^{1,8}$ of $K_{5,5}$ and consider the cycle $C = [1, 6, 2, 7, 3, 8, 4, 9, 5, 10]$ outlined by the full line. The generating set of C is $\mathcal{B}_C^{1,8} = \{C_{4,9}, C_{3,7}, C_{5,10}, C_{2,6}, C_{2,7}, C_{5,9}\}$, where $C_{xy} = [1, 8, x, y]$. $C_{3,7}$ is shown as dashed elementary cycle. One easily checks that $C \oplus C''$ is not an elementary cycle for each $C'' \in \mathcal{B}_C$. Therefore, there is no well-ordering of \mathcal{B}_C , and consequently $\mathcal{B}^{1,8}$ is not robust.

We can extend this negative result to all complete bipartite graphs $K_{m,n}$ with $m, n \geq 5$. By the same symmetry argument, it is always sufficient to consider the Kainen basis $\mathcal{K} = \mathcal{B}^{0,0'}$ with arbitrary fixed $0 \in V_1 = \{0, 1, \dots, m-1\}$ and $0' \in V_2 = \{0', 1', \dots, (n-1)'\}$. Furthermore, $K_{5,5}$ is an induced subgraph of $K_{m,n}$ for all $m, n \geq 5$. If 0 and $0'$ are vertices of $K_{5,5} \subset K_{m,n}$, then the Kainen basis of $K_{m,n}$ contains that of the induced and $K_{5,5}$. The elementary cycle C of Fig. 2 is also an elementary cycle in $K_{m,n}$ and its basis decomposition \mathcal{B}_C w.r.t. the Kainen basis of $K_{m,n}$ consists of the same basis elementary cycles as on $K_{5,5}$. Hence \mathcal{B}_C cannot be well-arranged on $K_{m,n}$ for $n, m \geq 5$.

Despite the counterexample above, we will show below that Kainen's assertion is true for small graphs.

Lemma 4. *The Kainen basis of $K_{m,n}$ with $m \leq 4$ and $n \leq 5$ is strictly robust.*

Proof. There is nothing to show for $n, m \leq 2$ since $K_{1,1}$ and $K_{1,2}$ have no cycles and $K_{2,2} = C_4$ consist of a single cycle. The longest cycles in $K_{m,n}$ has length $2 \min\{m, n\}$, i.e, we only need to consider cycles of length 4, 6, and 8. In the following, we show that for each type of cycle there is a well-arranged order of their Kainen basis elements. The correctness can be checked by direct computation in each case.

Cycles of length 4.

1. C contains $00'$. Then $C \in \mathcal{K}$ and there is nothing to show.
2. C contains exactly one of the vertices 0 or $0'$. W.l.o.g., $0 \in C$. Then $C = [0, i', j, j'] = C_{j,i'} \oplus C_{j,j'}$ and hence $\mathcal{B}_C = \{C_{j,i'}, C_{j,j'}\}$ is strictly well-arranged.
3. C contains neither 0 nor $0'$. All such elementary cycles are of the form $C = [i, i', j, j']$. Its strictly well-arranged sequence is $\mathcal{S}_C = (C_{j,i'}, C_{j,j'}, C_{i,j'}, C_{i,i'}, C)$.

Cycles of length 6.

1. C contains the edge $00'$. Such elementary cycles have the form $C = [0, 0', i, i', j, j']$. Its strictly well-arranged sequence is $\mathcal{S}_C = (C_{j,j'}, C_{j,i'}, C_{i,i'}, C)$.
2. C includes the ‘‘fixed vertices’’ 0 and $0'$ but not the edge $00'$. Such elementary cycles have the form $C = [0, i', i, 0', j, j']$. Its strictly well-arranged sequence is $\mathcal{S}_C = (C_{i,i'}, C_{j,j'}, C)$.
3. C contains exactly one of the vertices 0 or $0'$. W.l.o.g., $0 \in C$. Such elementary cycles have the form $C = [0, i', i, j', j, k']$. Its strictly well-arranged sequence is $\mathcal{S}_C = (C_{j,k'}, C_{j,j'}, C_{i,j'}, C_{i,i'}, C)$.
4. C contains neither 0 nor $0'$. Such elementary cycles have the form $C = [i, i', j, j', k, k']$. Its strictly well-arranged sequence is $\mathcal{S}_C = (C_{j,i'}, C_{j,j'}, C_{k,j'}, C_{k,k'}, C_{i,k'}, C_{i,i'}, C)$.

Cycles of length 8.

In our treatment these elementary cycles can only occur in $K_{4,4}$ and $K_{4,5}$. The case that C contains neither 0 nor $0'$ therefore does not appear.

1. C contains $00'$. All such elementary cycles C are of the form $C = [0, 0', i, i', j, j', k, k']$. Its strictly well-arranged sequence is $\mathcal{S}_C = (C_{j,i'}, C_{j,j'}, C_{k,j'}, C_{k,k'}, C_{i,i'}, C)$.
2. C includes $0 \in V_1$ and $0' \in V_2$ but not the edge $00'$. In this case we have two distinguish two kinds of elementary cycles: $C = [0, i', i, 0', j, j', k, k']$ or $D = [0, i', i, j, j', 0', k, k']$. The corresponding respective strictly well-arranged sequences are $\mathcal{S}_C = (C_{k,k'}, C_{k,j'}, C_{j,j'}, C_{i,i'}, C)$ and $\mathcal{S}_D = (C_{i,i'}, C_{i,j'}, C_{j,j'}, C_{k,k'}, D)$.
3. C contains exactly one of the vertices 0 or $0'$. W.l.o.g., $0 \in C$. All such elementary cycles are of the form $C = [0, i', i, j', j, k', k, l']$. Its strictly well-arranged sequence is $\mathcal{S}_C = (C_{k,l'}, C_{k,k'}, C_{j,k'}, C_{j,j'}, C_{i,j'}, C_{i,i'}, C)$.

□

4.2 The Cartesian Product $K_2 \square G$

In this subsection we will disprove a conjecture for cycle bases of Cartesian product graphs. The vertex set of the Cartesian product $G \square H$ of two graphs G and H is the set $\{(g, h) \mid g \in V(G), h \in V(H)\}$ that is the Cartesian product of the vertex sets of the factors. Two vertices $(g_1, h_1), (g_2, h_2)$ are adjacent in $G \square H$ if $g_1 g_2 \in E(G)$ and $h_1 = h_2$ or if $g_1 = g_2$ and $h_1 h_2 \in E(H)$. For more detailed informations we refer the interested reader to [8].

In [9] it is conjectured that robust cycle bases can be constructed for Cartesian product graphs of the form $T \square G$, where T is a tree, provided a robust basis \mathcal{R} is already known for G . More precisely, the basis \mathcal{R}' consisting of the known robust basis \mathcal{R} for a single G -fiber together with all quadrangles of the form $e \square f$, $e \in E(T)$, $f \in E(G)$ has been proposed.

In the following we consider the planar graph G shown in Figure 3.

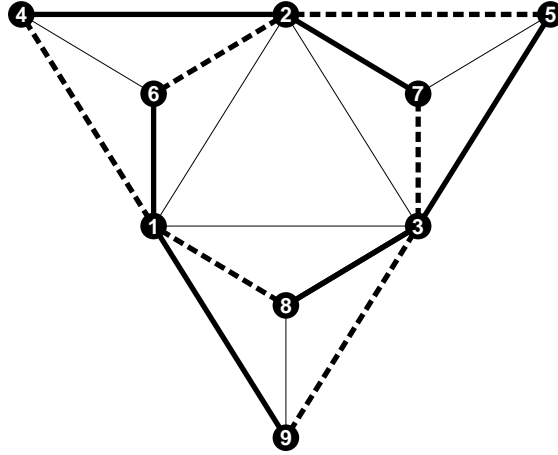


Figure 3: Planar graph G . The Cartesian product $K_2 \square G$ consists of two copies of G in which corresponding vertices are connected by an edge. We consider the elementary cycle C which is partially contained in both the upper copy of G (thick edges) and in the lower copy of G (dashed edges). The elementary cycle C changes from upper to lower edges along a K_2 -fiber wherever a *single* dashed and a *single* thick edge meet, i.e. not on vertices 1, 2 and 3.

The Cartesian product $K_2 \square G$ has 18 vertices and 45 edges, i.e., $\mu(K_2 \square G) = 28$. The basis \mathcal{R}' consists of the faces of the facial basis and all quadrangles of the form $K_2 \square f$, $f \in E(G)$. As mentioned in the introduction, any facial basis of a planar graph is robust [5], and hence also quasi-robust.

Now consider the elementary cycle C indicated in Fig. 3. For each triangle $D \in \mathcal{R}$, $D \oplus C$ is not an elementary cycle because we always obtain at least one vertex of degree 4. The same is true for any quadrangle $Q = K_2 \times f$ with $f \in G$. Again we obtain at least one vertex with degree larger than 2 in all cases. Thus the elementary cycle C cannot be constructed from any elementary cycle in $K_2 \square G$ by adding an element of the basis \mathcal{R}' . Thus the cycle basis \mathcal{R}' is not quasi-robust.

As robust implies quasi-robust, \mathcal{R}' is in particular also not a robust basis, disproving Kainen's conjecture.

5 Quasi-Robust Basis for $K_{m,n}$

Theorem 5. *The Kainen basis \mathcal{K} of $K_{m,n}$ is quasi-robust for all m, n .*

Proof We first show that each cycle C in $K_{m,n}$ can be obtained from an elementary cycle that includes the “fixed edge” $00'$ by adding a single basis element $D \in \mathcal{K}$.

Case 1: C includes the “fixed vertices” $0 \in V_1$ and $0' \in V_2$ but not the edge $00'$.

Claim. There is a $D \in \mathcal{K}$ such that $C \oplus D$ is an elementary cycle containing $00'$ satisfying $|C \oplus D| = |C|$.

Proof. We choose a cycle C through 0 and $0'$ that does not include the edge $00'$. W.l.o.g. we can label the vertices as such that $C = [0, 1', 1, \dots, r', r, 0', r+1, (r+1)', \dots, s, s']$. Now we choose $D = [0, 0', r+1, 1] \in \mathcal{K}$ and compute $C \oplus D = [0, 0', r, r', \dots, 2', 1, 1', r+1, (r+1)', \dots, s, s', 0]$, which we can also interpret as the closed path $(0, 0', r, r', P)$ where P is a path in $K_{p,q}$ which (1) connects 0 and r' , (2) does not include $0', r$, and (3) has maximum vertex degree 2. Thus $C \oplus D$ is an elementary cycle. We have added the two edges $00'$ and $1'(r+1)$ and removed the edges $01'$ and $0'(r+1)$, hence the length of the cycle C is preserved. \square

Case 2: C contain exactly one of the vertices 0 or $0'$.

Claim. There is a $D \in \mathcal{K}$ such that $C \oplus D$ is an elementary cycle containing the edge $00'$ satisfying $|C \oplus D| = |C|$.

Proof. W.l.o.g. we assume that C contains 0 , i.e., $[0, i', k, P, j']$, where P is a path connecting k and j' . Set $D = [0, 0', k, i']$. Then $C \oplus D$ is an elementary cycle of length $|C|$ that contains $00'$. \square

Case 3: C contains neither 0 nor $0'$.

Claim. There is a $D \in \mathcal{K}$ such that $C \oplus D$ is an elementary cycle containing $00'$ of length $|C| + 2$.

Proof. Let i and i' be any two adjacent vertices of C . The $D = [0, 0', i, i'] \in \mathcal{K}$ and $C \oplus D$ is again an elementary cycle. Since we add three edge and delete 1, its length is $|C| + 2$. \square

Now we show that each elementary cycle through the edge $00'$ can be obtained in two steps from a strictly shorter elementary cycle with the same property.

Claim. Let C be an elementary cycle of length $|C| \geq 6$ including the edge $00'$. Then there are basis cycles $D, E \in \mathcal{K}$ such that $C \oplus D$ is an elementary cycle of length $|C \oplus D| = |C| - 2$ and $C \oplus D \oplus E$ is an elementary cycle of length $|C \oplus D \oplus E| = |C| - 2$ that includes the edge $00'$.

Proof. We can write C in the form $C = [0, 0', j, j', P]$, where P is a path of length $|P| \geq 2$ in $K_{m,n}$ that (1) connects 0 and j' , (2) does not include the edge $0'j$, and (3) has maximum vertex degree two. Choosing $D = [0, 0', j, j'] \in \mathcal{K}$ we immediately compute $C \oplus D = [0, j', P]$. Obviously, D is an elementary cycle of length $|C| - 2$. Since P has length at least 2, we can re-write this in the form $C \oplus D = [0, j', P] = [0, j', k, k', Q]$, where Q is a path of length $|P| - 2 \geq 0$ connecting k' and 0 . Now we choose a second basis cycle $E = [0, 0', k, j'] \in \mathcal{K}$ and obtain $C \oplus D \oplus E = [0, 0', k, k', Q]$. This is again an elementary cycle, contains the edge $00'$ and has length $|C| - 2$. \square

Starting from an arbitrary elementary cycle in $K_{m,n}$, we obtain in at most one step a cycle that runs through $00'$. Then we can reduce its length by two in a pair of steps, thereby arriving at a cycle

of length 4 through $00'$ in not more than $|C| - 1$ operations. Since all quadrangles containing $00'$ are by construction elements of \mathcal{K} , we can reduce every elementary cycle C to \emptyset in no more than $|C|$ steps in such a way that all intermediates are also elementary cycles. Reversing the order, we have thus constructed a well-arranged sub-sequence $\mathcal{S}(C)$ of elements of \mathcal{K} that generates C . **Q.E.D.**

Lemma 6. *The Kainen basis \mathcal{K} of $K_{m,n}$ is not strictly quasi-robust for all $m, n \geq 5$.*

Proof. In order to construct a counterexample, we again choose a cycle C of length 10 through 0 and $0'$ that does not include the edge $00'$. W.l.o.g. we can label the vertices as such that $C = [0, 1', 3, 3', 1, 0', 2, 4', 4, 2']$. From Theorem 5 follows that there is a well arranged sequence \mathcal{S}_C of the Kainen basis elements. Let $\mathcal{S}_C = (\dots, D, C)$ be any such sequence for the particular chosen elementary cycle C . First notice that any second last element $D \in \mathcal{S}_C$ must include the edges $0j'$ and $0'i$, both included in C , otherwise $C \oplus D$ would not be an elementary cycle and thus $\mathcal{S}_C = (\dots, D, C)$ would not be well-arranged. Therefore we have exactly four cases: $D_{1,1'}$, $D_{2,2'}$, $D_{1,2'}$ and $D_{2,1'}$. By symmetry we can reduce these cases to $D_{i,i'}$ and $D_{i,j'}$.

First notice that $D_{i,i'}$ cannot be the second last element of \mathcal{S}_C , since $C \oplus D_{i,i'}$ crumbles down into two elementary cycles.

Thus we only have to consider the case $D_{i,j'}$. By direct computing we get the elementary cycle $Q = C \oplus D_{i,j'}$ and thus we can observe that $D_{i,j'}$ is a possibly proper element of \mathcal{S}_C . But on the other hand we can observe that $Q \cap D_{i,j'} = \{\{0, 0'\}, \{i, j'\}\}$ consisting of two paths. Thus none of the well-arranged sequences \mathcal{S}_C can be strictly well-arranged, from what the assumption follows. \square

6 Concluding Remarks

A major obstacle in the analysis of robust cycle bases is that no generally applicable “tools” are known, a difficulty that is confounded by observation that robustness properties are unrelated to other, better-studied, properties of cycle bases such as fundamentality or minimality [11]. By relaxing the definition to quasi-robustness we abandon the requirement that a cycle must be constructed from the basis elements in its linear representation only. This buys us a generic strategy for constructive proofs to establish quasi-robust: It suffices to show that every elementary cycle can be reduced by \oplus -addition of basis elements to another elementary cycle C' that is “closer to” \emptyset w.r.t. to some (partial) ordering of elementary cycles (in our example simply by length). As an example, we directly prove the following corollary of the robustness result in [5]:

Lemma 7. *The facial basis of planar graph G is quasi-robust.*

Proof. All elementary cycles of G are closed intersection-free curves in the plane. Define a partial order on the elementary cycles such that $C \prec C'$ if $C \neq C'$ lies inside of C' , and set $\emptyset \prec C$ for all elementary cycles. Given C , there is a planar face F such that $C \oplus F \prec C$ is again an elementary cycle. Since there are only finitely many elementary cycles, every elementary cycle can be reduced to \emptyset . \square

In addition to this technical simplification, quasi-robustness is of interest in its own right in the context of sampling algorithms on the cycle space: quasi-robustness is necessary and sufficient for ergodicity of the Markov chains considered in [10].

We have briefly touched upon the “strict” variants of robustness and quasi-robustness. It stands to reason that the additional geometric constraint — intersection of elementary cycles with a basis a cycle in a single path — further complicates constructive proofs. Again, strict quasi-robustness seems much easier to handle than strict robustness.

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